

On the computation of fusion over the affine Temperley-Lieb algebra

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intro:

fusion for finite associative algebras

operator product expansion and fusion

- operator product expansion in CFT seen as a relation between primary or quasi-primary fields

operator product expansion and fusion

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- formal product between representations of conformal algebra (Virasoro, affine, ...)

operator product expansion and fusion

- operator product expansion in CFT seen as a relation between primary or quasi-primary fields
- formal product between representations of conformal algebra (Virasoro, affine, ...)
- properties that are verified or might be desired:
 1. **binary relation:** $M \times_f N \in \text{mod } A$ for $M, N \in \text{mod } A$
 2. **commutativity:** $M \times_f N \simeq N \times_f M$
 3. **associativity:** $(M \times_f N) \times_f O \simeq M \times_f (N \times_f O)$
 4. **finitess:** $M \times_f N$ is a finite sum of indecomposables

fusion for finite 2d models

- 2d lattice models: percolation, Ising, dimer, loop models,
 H_{XXZ}

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- e.g.: $H_{XXZ}(n) : (\mathbb{C}^2)^{\otimes n} \rightarrow (\mathbb{C}^2)^{\otimes n}$ where $H_{XXZ}(n) = -\sum_{i < n} \rho(e_i)$ (where $e_i, 1 \leq i < n$, generate $TL_n(q + q^{-1})$) with

$$\rho(e_i) = \frac{1}{2} \left(\sigma_i^x \sigma_{i+1}^x + \sigma_i^y \sigma_{i+1}^y - \frac{1}{2}(q + q^{-1})(\sigma_i^z \sigma_{i+1}^z - I) + \frac{1}{2}(q - q^{-1})(\sigma_i^z - \sigma_{i+1}^z) \right)$$

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- underlying algebras in many models: Temperley-Lieb algebra(s)

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how to define fusion of modules over TL_n ?

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how to define fusion of modules over TL_n ?

- “natural product”: tensor product of representations
- works for groups, Lie algebras, Hopf algebras (through the coproduct), ...
- has all the desired properties (binary operation, commutativity, associativity, finiteness), at least under reasonable hypotheses;
- ... **but fails for TL_n !**

Read & Saleur's idea based on $TL_m \times TL_n \subset TL_{m+n}$

Read & Saleur (2007), Gainutdinov & Vasseur (2013)

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For $M \subset \text{mod } TL_m, N \subset \text{mod } TL_n$ and $(M, N) \subset \text{mod } TL_m \times TL_n$,
define

$$\begin{aligned} M \times_f N &\stackrel{\text{def}}{=} (M, N) \uparrow_{TL_m \times TL_n}^{TL_{m+n}} \\ &= TL_{m+n} \otimes_{TL_m \times TL_n} (M \otimes_{\mathbb{C}} N) \in \text{mod } TL_{m+n} \end{aligned}$$

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the fusion \times_f satisfies

- commutativity
- associativity
- finiteness

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the fusion \times_f satisfies

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but

- it is not a binary relation on a given $\text{mod } TL_n$

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the property of *binary relation* should be replaced by *stability*:

the fusion of standard modules S_{m,k_1} and S_{n,k_2} depends only on k_1 and k_2 when m or n are large enough

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(compare with: the limit $n \rightarrow \infty$ of the spectrum of $H_{XXZ}(n)$ restricted to $S_{n,k}$ exists)

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Theorem (Belletête (2015))

The fusion of irreducible $TL_{\bullet}(q + q^{-1})$ -modules $I_{m,i}$ and $I_{n,j}$ with $0 \leq i, j < \ell$ and $q = e^{\pi i / \ell}$ reproduces the fusion of primary fields in the minimal models $\mathcal{M}(\ell + 1, \ell)$ under the identification $I_{m,i} \rightarrow \phi_{1,1+i}$.

regular and affine TL algebras, and their modules

the regular TL_n

$$TL_n(\beta = q + q^{-1}) = \\ \langle \text{id}, e_i, 1 \leq i < n \rangle$$

$$e_i^2 = \beta e_i$$

$$e_i e_{i \pm 1} e_i = e_i$$

$$e_i e_j = e_j e_i, \quad |i - j| > 1$$

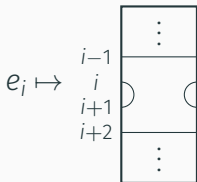
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an example for $n = 10$

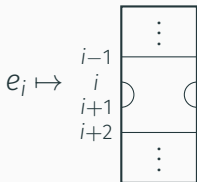
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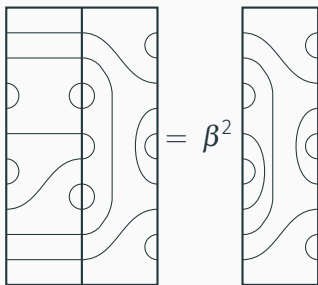
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multiplication in $TL_{10}(\beta)$



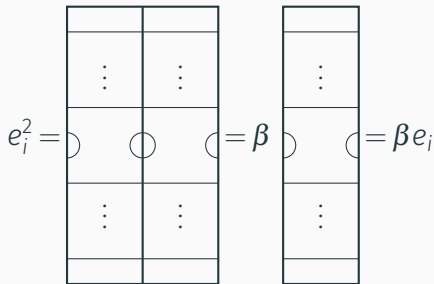
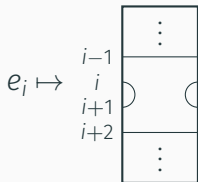
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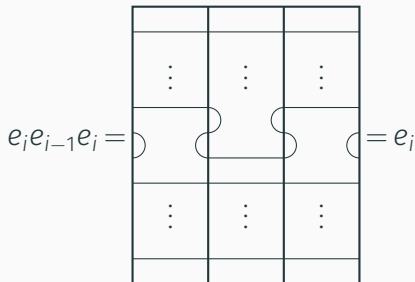
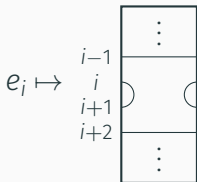
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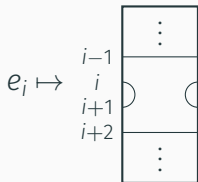
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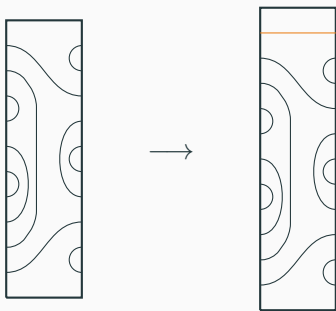
TL_n as a diagram algebra

$TL_n \simeq$ formal linear combinations of n -diagrams (n links joining $2n$ points) with above multiplication.

$$\begin{aligned} \dim TL_n &= \frac{1}{n+1} \binom{2n}{n} \\ &= n\text{-th Catalan number} \end{aligned}$$

natural inclusion $TL_n \subset TL_{n+1}$

$\iota : TL_n \rightarrow TL_{n+1}$ by adding a through line on top :



the affine TL_n^a

$$TL_n^a(\beta = q + q^{-1}) = \langle \text{id}, \tau, \tau^{-1}, e_i, 1 \leq i \leq n \rangle$$

$$e_i^2 = \beta e_i, \quad e_i e_{i \pm 1} e_i = e_i,$$

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$$\tau e_i = e_{i+1} \tau, \quad e_1 \tau^2 = e_1 e_2 \dots e_{n-1}$$

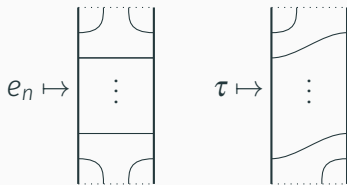
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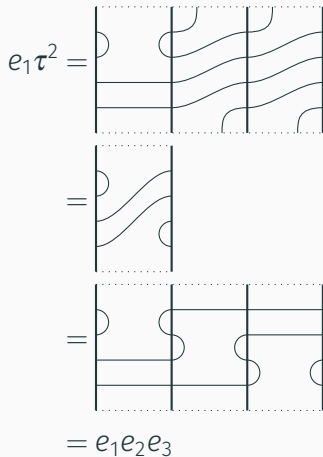
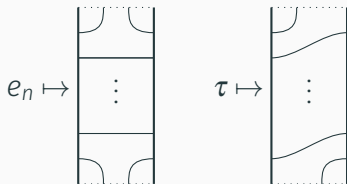
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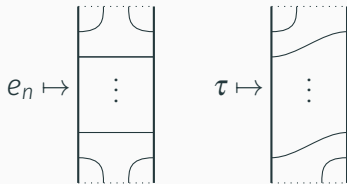
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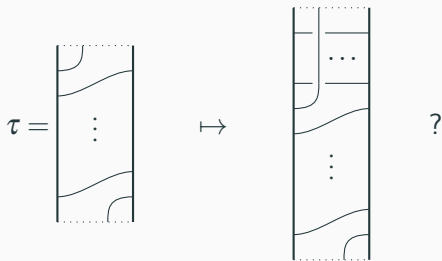
TL_n^a as a diagram algebra

$TL_n^a \simeq$ formal linear combinations of affine n -diagrams with concatenation as multiplication.

$$\dim TL_n^a = \infty$$

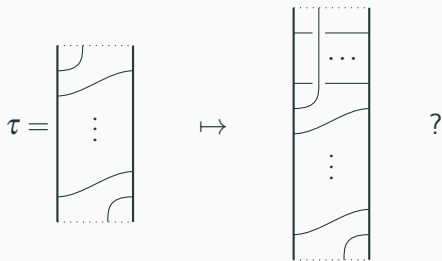
is $TL_m^a \subset TL_n^a$ when $m \leq n$?

- TL_m^a is not diagrammatically included into TL_n^a :



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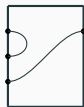


- Gainutdinov & Saleur (2016) have found an injective morphism $TL_m^a \times TL_n^a \rightarrow TL_{m+n}^a$ and were able to propose a fusion product along the line of \times_f ;
- our fusion products are introduced through another path.

modules over TL_n

(n, k) -diagrams have n nodes on the left, k on the right.

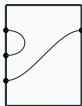
Here is a $(3, 1)$ -diagram:



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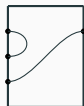
Standard module $S_{n,k}, n \geq k$

Formal linear combinations of (n,k) -diagrams with k through lines.

modules over TL_n

(n,k) -diagrams have n nodes on the left, k on the right.

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action is by concatenation

$$\begin{array}{|c|} \hline \bullet \\ \hline \bullet \\ \hline \bullet \\ \hline \end{array} \times \begin{array}{|c|} \hline \bullet \\ \hline \\ \hline \\ \hline \end{array} = \begin{array}{|c|c|} \hline \bullet & \bullet \\ \hline & \\ \hline & \\ \hline & \\ \hline \end{array} = \beta \begin{array}{|c|} \hline \bullet \\ \hline \\ \hline \\ \hline \end{array}$$

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but

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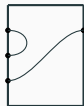
Standard module $S_{n,k}, n \geq k$

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modules over TL_n

(n,k) -diagrams have n nodes action is by concatenation on the left, k on the right.

Here is a $(3,1)$ -diagram:



$$\begin{array}{|c|} \hline \bullet \\ \hline \bullet \\ \hline \bullet \\ \hline \end{array} \times \begin{array}{|c|} \hline \bullet \\ \hline \bullet \\ \hline \bullet \\ \hline \end{array} = \begin{array}{|c|c|} \hline \bullet & \bullet \\ \hline \bullet & \bullet \\ \hline \bullet & \bullet \\ \hline \end{array} = \beta \begin{array}{|c|} \hline \bullet \\ \hline \bullet \\ \hline \bullet \\ \hline \end{array}$$

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Standard module $S_{n,k}, n \geq k$

Formal linear combinations of (n,k) -diagrams with k through lines.

Irreducible module $I_{n,k}$

Quotient of $S_{n,k}$ by its radical (quotient is irreducible).

I. modules over TL_n^a

Affine standard module $S_{n,k}^a$

Formal linear combinations of affine (n, k) -diagrams.

I. modules over TL_n^a

Affine standard module $S_{n,k}^a$

Formal linear combinations of affine (n, k) -diagrams.

$$\left\{ \dots, v_1 \tau = \begin{array}{|c|} \hline \text{Diagram 1} \\ \hline \end{array}, v_1 = \begin{array}{|c|} \hline \text{Diagram 2} \\ \hline \end{array}, v_1 \tau^{-1} = \begin{array}{|c|} \hline \text{Diagram 3} \\ \hline \end{array}, \dots, \right.$$
$$\dots, v_2 \tau = \begin{array}{|c|} \hline \text{Diagram 4} \\ \hline \end{array}, v_2 = \begin{array}{|c|} \hline \text{Diagram 5} \\ \hline \end{array}, v_2 \tau^{-1} = \begin{array}{|c|} \hline \text{Diagram 6} \\ \hline \end{array}, \dots,$$
$$\dots, v_3 \tau = \begin{array}{|c|} \hline \text{Diagram 7} \\ \hline \end{array}, v_3 = \begin{array}{|c|} \hline \text{Diagram 8} \\ \hline \end{array}, v_3 \tau^{-1} = \begin{array}{|c|} \hline \text{Diagram 9} \\ \hline \end{array}, \dots,$$
$$\dots, v_4 \tau = \begin{array}{|c|} \hline \text{Diagram 10} \\ \hline \end{array}, v_4 = \begin{array}{|c|} \hline \text{Diagram 11} \\ \hline \end{array}, v_4 \tau^{-1} = \begin{array}{|c|} \hline \text{Diagram 12} \\ \hline \end{array}, \dots \left. \right\}$$

II. modules over TL_n^a : the cell modules (Graham & Lehrer)

morphisms on $S_{n,k}^a$

For $z \in \mathbb{C}^\times$, define

$$\begin{aligned} f_{n,k}^z : S_{n,k}^a &\longrightarrow S_{n,k}^a, & f_{n,0}^z : S_{n,0}^a &\longrightarrow S_{n,0}^a \\ &: v \mapsto v \circ (\tau - z \text{id}), & &: v \mapsto v \circ (o - (z + z^{-1}) \text{id}) \\ &(\text{for } k > 0) & &(\text{for } k = 0) \end{aligned}$$

where o is the map that adds a non-contractible loop.

II. modules over TL_{-n}^a : the cell modules (Graham & Lehrer)

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cell modules $S_{n,k;z}^a$

The cell modules $S_{n,k;z}^a$ are the quotients of $S_{n,k}^a / \ker f_{n,k}^z$.

$$\dim S_{n,k;z}^a = \binom{n}{(n-k)/2}.$$

functors and fusions

braiding on $\tilde{\mathcal{T}}\mathcal{L}$

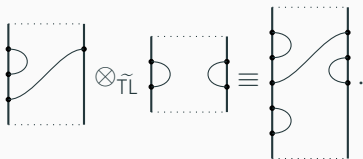
$\tilde{\mathcal{T}}\mathcal{L}$

the set of (m, r) -diagrams,
 $m, r \in \mathbb{N}$, with
concatenation \times and
superposition \otimes

braiding on $\tilde{\mathbb{L}}$

$\tilde{\mathbb{L}}$

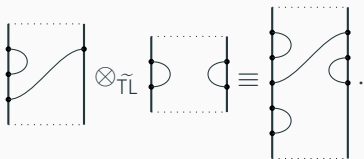
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braiding is a natural isomorphism

$$\eta : -_1 \otimes -_2 \longrightarrow -_2 \otimes -_1$$

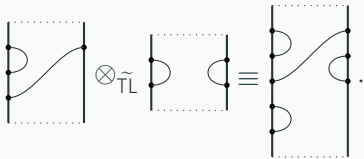
It braids as follows: for f
a (m, r) – diagram and g a
 (n, s) – diagram, then

$$\eta_{m,n}(f \otimes g) = (g \otimes f)\eta_{r,s}$$

braiding on $\widetilde{\text{TL}}$

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the set of (m, r) -diagrams,
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braiding is a natural isomorphism

$$\eta : -1 \otimes -2 \longrightarrow -2 \otimes -1$$

It braids as follows: for f a (m, r) -diagram and g a (n, s) -diagram, then

$$\eta_{m,n}(f \otimes g) = (g \otimes f)\eta_{r,s}$$

(actual form not necessary;
based on WL Chow (1948);
used by P Martin (1991) for
TL.)

maps $TL_n \rightarrow TL_n^a$ and $TL_n^a \rightarrow TL_n$

TL_n is a subalgebra of TL_n^a .

Theorem

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Define $\phi : TL_n^a \rightarrow TL_n$ by

$$\phi(\tau) = (-q)^{3/2} \eta_{n-1,1},$$

$$\phi(e_i) = e_i, \quad i < n,$$

$$\phi(e_n) = \eta_{n-1,1}^{-1} e_1 \eta_{n-1,1}$$

if $n \geq 2$; for $n = 2$, by

$$\phi(\tau) = (-q)^{3/2} \eta_{1,1}, \quad \phi(e_2) = 0;$$

and for $n = 1$, by

$$\phi(\tau) = (-q)^{3/2} 1_{TL_1}.$$

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The map $\phi : TL_n^a \rightarrow TL_n$ is a surjective morphism of algebras.

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induction and restriction

The inclusion ι gives two functors:

$$\uparrow_r^a : \text{mod TL}_n \rightarrow \text{mod TL}_n^a, \quad \downarrow_r^a : \text{mod TL}_n^a \rightarrow \text{mod TL}_n.$$

induction and restriction

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The surjective morphism $\phi : \text{TL}_n^a \rightarrow \text{TL}_n$ gives two more:

$$\uparrow_a^r : \text{mod } \text{TL}_n^a \rightarrow \text{mod } \text{TL}_n, \quad \downarrow_a^r : \text{mod } \text{TL}_n \rightarrow \text{mod } \text{TL}_n^a.$$

The latter are

$$V \mapsto \uparrow_a^r(V) = {}_{\text{TL}_n}(\text{TL}_n)_{\text{TL}_n^a} \otimes_{\text{TL}_n^a} V, \quad \text{for } V \in \text{mod } \text{TL}_n^a$$

and

$$X \mapsto \downarrow_a^r(X) = \text{Hom}_{\text{TL}_n}(\text{TL}_n, X), \quad \text{for } X \in \text{mod } \text{TL}_n.$$

an example: computing $\uparrow_a^r S_{3,1}^a$

The module $\uparrow_a^r S_{3,1}^a$ is

$$\mathrm{TL}_3 \otimes_{\mathrm{TL}_3^a} S_{3,1}^a = \left\{ a \otimes_{\mathrm{TL}_3^a} v \mid a \in \mathrm{TL}_3 \text{ and } v \in S_{3,1}^a \right\}$$

where

$$a \otimes_{\mathrm{TL}_3^a} bv = a\phi(b) \otimes_{\mathrm{TL}_3^a} v, \quad \text{for any } b \in \mathrm{TL}_3^a.$$

A basis for $S_{3,1}^a$ is

$$\left\{ \begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} \tau^i, \begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \end{array} \tau^i, i \in \mathbb{Z} \right\}.$$

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But

$$\boxed{\tau} = \boxed{\text{diagram}} = \boxed{\text{diagram}} = \boxed{\text{diagram}} = \tau^3 \boxed{\tau}$$

and thus

$$a \otimes_{\mathrm{TL}_3^a} \boxed{\tau^j} = a \otimes_{\mathrm{TL}_3^a} \tau^{3i} \boxed{\tau} = \underbrace{a \phi(\tau^{3i})}_{\in \mathrm{TL}_3} \otimes_{\mathrm{TL}_3^a} \boxed{\tau}.$$

an example: computing $\uparrow_a^r S_{3,1}^a$

The module $\uparrow_a^r S_{3,1}^a$ is

$$TL_3 \otimes_{TL_3^a} S_{3,1}^a = \left\{ a \otimes_{TL_3^a} v \mid a \in TL_3 \text{ and } v \in S_{3,1}^a \right\}$$

Finally

$$\uparrow_a^r S_{3,1}^a = \text{span} \left\{ \text{id} \otimes_{TL_3^a} \begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \square \\ \hline \end{array}, e_2 \otimes_{TL_3^a} \begin{array}{|c|} \hline \square \\ \hline \square \\ \hline \square \\ \hline \end{array} \right\}$$

and

$$\uparrow_a^r S_{3,1}^a = S_{3,1}.$$

fusion products

\times_f^1 and \times_f^2

For $M \in \text{mod TL}_m^a$ and $N \in \text{mod TL}_n^a$, define

$$M \times_f^1 N \stackrel{\text{def}}{=} \Downarrow_a^r ((\Uparrow_a^r M) \times_f (\Uparrow_a^r N)),$$

$$M \times_f^2 N \stackrel{\text{def}}{=} \Uparrow_r^a ((\Uparrow_a^r M) \times_f (\Uparrow_a^r N)).$$

fusion products

Because of

Proposition

The functors \uparrow_a^r , \uparrow_r^a and \downarrow_a^r satisfy

$$\uparrow_a^r \circ \uparrow_r^a \xrightarrow{\sim} \uparrow_a^r \circ \downarrow_a^r \xrightarrow{\sim} \text{id}_{\text{mod TL}_n} \quad \text{and, if } n \neq 2, \quad \downarrow_r^a \circ \downarrow_a^r \xrightarrow{\sim} \text{id}_{\text{mod TL}_n}.$$

these fusion products are

- commutative;
- associative.

Do they satisfy **finiteness** and **stability**?

computing affine fusion

\downarrow_a^r and \uparrow_a^r : on standards and cells

Proposition

If $n \neq 2$,

$$\uparrow_a^r S_{n,k}^a \simeq S_{n,k},$$

and if $n = 2$,

$$\uparrow_a^r S_{2,2}^a \simeq M_{2,2}, \quad \uparrow_a^r S_{2,0}^a \simeq \begin{cases} P_{2,2}, & \text{if } \beta = 0, \\ 0, & \text{otherwise.} \end{cases}$$

\downarrow_a^r and \uparrow_a^r : on standards and cells

Proposition

$$\uparrow_a^r S_{n,k}^a \simeq S_{n,k}, \quad \text{if } n \neq 2.$$

\downarrow_a^r and \uparrow_a^r : on standards and cells

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Proposition

$$\uparrow_a^r S_{n,k;z}^a \simeq \begin{cases} S_{n,k}, & \text{if } z = (-q)^{\frac{k+2}{2}}, \\ 0, & \text{otherwise,} \end{cases} \quad \text{if } n \neq 2.$$

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Induction \uparrow_a^r of an affine standard is always trivial **unless** the two parameters z and q are perfectly tuned.

\downarrow_a^r and \uparrow_a^r : on irreducibles

We shall use the notation

$$z_k = (-q)^{\frac{k}{2}}.$$

\Downarrow_a^r and \Uparrow_a^r : on irreducibles

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Proposition

$$\Downarrow_a^r I_{n,k} \simeq I_{n,k;z_{k+2}}^a \quad \text{and} \quad I_{n,k} \simeq \Uparrow_a^r I_{n,k;z_{k+2}}^a, \quad \text{if } n \neq 2.$$

\downarrow_a^r and \uparrow_a^r : on irreducibles

We shall use the notation

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$$\downarrow_a^r I_{n,k} \simeq I_{n,k; z_{k+2}}^a \quad \text{and} \quad I_{n,k} \simeq \uparrow_a^r I_{n,k; z_{k+2}}^a, \quad \text{if } n \neq 2.$$

Proposition

$$\downarrow_a^r S_{n,k} \simeq S_{n,k; z_{k+2}}^a / S_{n,k+2; z_k}^a, \quad \text{if } n \neq 2.$$

\uparrow_r^a on regular standards

Proposition

(i) Let $\beta \neq 0$. For $k \geq 2$, the short sequence

$$0 \longrightarrow \uparrow_r^a S_{n,k-2} \xrightarrow{i_k} \uparrow_r^a S_{n,k} \longrightarrow S_{n,k}^a \longrightarrow 0$$

is exact and $\uparrow_r^a S_{n,k} \simeq S_{n,k}^a$ for $k = 0, 1$.

(ii) $\uparrow_r^a S_{n,k}$ is indecomposable.

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(ii) $\uparrow_r^a S_{n,k}$ is indecomposable.

Proposition (Peirce decomposition of TL_n^a)

For generic q

$$TL_n^a \simeq \bigoplus'_{0 \leq k \leq n} \dim(S_{n,k}) \cdot \uparrow_r^a S_{n,k}$$

as a left module.

Proposition

If q is generic and $0 < t \leq (n - k)/2$,

$$\downarrow_r^a S_{n,k;z}^a \simeq \bigoplus_{r=0}^{(n-k)/2} S_{n,k+2r}, \quad \downarrow_r^a I_{n,k;z_{k+2t}}^a \simeq \bigoplus_{r=0}^{t-1} S_{n,k+2r}.$$

\downarrow_r^a on affine standards

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If q is generic and $0 < t \leq (n - k)/2$,

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(no result yet for the cases when q is a root of unity)

examples of fusion

Proposition

Let q be generic, $r \geq s \geq 0$ and $n, m > 2$. Then

$$S_{n,r}^a \times_f^1 S_{m,s}^a \simeq I_{n,r;Z_{r+2}}^a \times_f^1 I_{m,s;Z_{s+2}}^a \simeq \bigoplus'_{t=r-s}^{r+s} I_{n+m,t;Z_{t+2}}^a,$$

$$S_{n,r}^a \times_f^2 S_{m,s}^a \simeq I_{n,r;Z_{r+2}}^a \times_f^2 I_{m,s;Z_{s+2}}^a \simeq \bigoplus'_{t=r-s}^{r+s} \uparrow_r^a S_{n+m,t},$$

$$S_{n,r;Z_{r+2}}^a \times_f^1 S_{m,s;Z_{s+2}}^a \simeq \bigoplus'_{t=r-s}^{r+s} I_{n+m,t;Z_{t+2}}^a,$$

$$I_{n,r;Z_{r+2a}}^a \times_f^3 I_{m,s;Z_{s+2b}}^a (m) \simeq \bigoplus_{x=0}^{a-1} \bigoplus_{y=0}^{b-1} (I_{n,r+2x;Z_{r+2(x+1)}}^a \times_f^1 I_{m,s+2y;Z_{s+2(y+1)}}^a)$$

with \bigoplus' meaning a direct sum with an increment of 2.

concluding remarks

done- and to-do-lists

done

- two definitions of affine fusion \times_f^1 and \times_f^2 (based on regular fusion \times_f);
- commutativity, associativity, finiteness, stability: ✓ ;
- computed for q generic.

to do

- compute fusion \times_f^1 and \times_f^2 when q is a root of unity (might need more representation theory of TL_n^a);
- explore other definitions. For example $M \times_f^3 N = \Downarrow_a^r ((\Downarrow_r^a M) \times_f (\Downarrow_r^a N))$;
- understand relations between \times_f^{GS} and $(\times_f^1$ or $\times_f^2)$;
- compare with known cases of fusion in CFT (start with modules arising in H_{XXZ} ?).